

# Soft-error Mitigation for Asynchronous FPGAs

Wonjin Jang, Alain J. Martin  
 Computer Science Department  
 California Institute of Technology  
 Pasadena, CA 91125

**Abstract**—This paper addresses the issue of soft errors in quasi delay-insensitive (QDI) FPGA. We combine two soft-error mitigation schemes. One is to duplicate and double-check computation cells, and the other is to interlock coupled inverters of programmable bits. We present a soft-error tolerant logic cells of QDI FPGAs based on the schemes.

## I. INTRODUCTION

As the circuit feature size decreases, soft errors become an important issue for designing logic cells and memories [1]. Quasi-delay-insensitive (QDI) circuits are the family of asynchronous circuits that operate with the weakest timing assumption (isochronic fork) [2]. While QDI circuits are quite robust to variable operating conditions, they are susceptible to soft errors like any other digital circuits. To exploit general robustness of QDI systems, we have developed a scheme to protect QDI circuits from soft errors by duplicating and double-checking of nodes [3].

FPGAs (Field Programmable Gate Arrays) become more attractive than before because of recent enhancement of capacity and performance. However, concern for clock distribution and increased power consumption of synchronous FPGAs is growing. The problems lead us to consider soft-error tolerant QDI FPGAs because of their clockless and power-efficient features. Some general QDI FPGA architectures have been proposed [4][5]. The purpose of this paper is to harden the Caltech FPAG architecture to obtain soft-error tolerance.

## II. DESIGN FOR SOFT-ERROR TOLERANT FPGA CELL

### A. FPGA Architecture

A basic FPGA tile consists of a cluster, two connection boxes (C-box) and a switch box (S-box). A system described in a high-level language is decomposed into implementable modules, which correspond to clusters, and whose interconnection information is mapped into C-boxes and S-boxes. C-boxes is for connecting a cluster to interconnect-paths and S-boxes is for switching interconnect-paths. Logic cells in a cluster share inputs and outputs, and each logic cell is a PCHB (Pre-charged Half Buffer) [6], which consists of a pulldown computation stack, validity trees of inputs and outputs, and so on. The cell contains programmable SRAM bits: some of the bits are for configuring computations and the others are for setting patterns of communication between cells.

There are two types of soft errors in FPGA: an error in computation parts (e.g., pulldown stacks) can bring about wrong computations temporarily, and an error in programmable bits

may change system configurations permanently as well as lead to wrong computations.

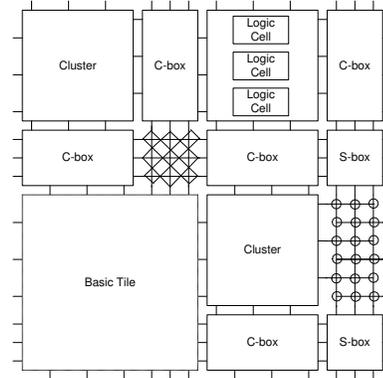


Fig. 1. FPGA Architecture

### B. Duplicated Double-checking Logic Cell

The key idea of detecting or correcting errors in any system is redundancy. Triple modular redundancy (TMR) with voters is widely used to correct an error in synchronous circuits. But in QDI circuits dual copies are enough to restore corrupt data due to the stability property of QDI circuits. The stability property is a system property that if the assignment of output of a gate starts, inputs of the gate hold current values until the assignment is completed.

Let us define a *duplicated double-checking* (DD) circuit. To get a DD circuit, we duplicate all gates in the original circuit and double-check all output nodes. Double-checking duplicated output nodes  $z_a, z_b$  means that we replace  $z_a, z_b$  with new nodes (e.g.,  $z'_a, z'_b$ ) and introduce two C-elements that share the inputs  $z'_a, z'_b$  and whose outputs are  $z_a$  and  $z_b$ . Figure 2 shows what a DD gate looks like. DD circuits can

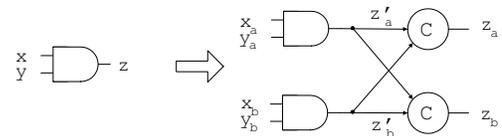


Fig. 2. AND gate and DD AND Gate

tolerate multiple soft errors unless two or more errors occur in two adjacent DD gates[7].

The soft-error-tolerance of DD circuits is based on the fact that at least one duplicated variable of each pair in DD

circuits will contain a correct value, and the double-checking scheme prevents corrupted values from propagating to subsequent gates. That is, an error is confined between double-checking C-elements, and the correct data can be reconstructed due to the stability property. Although soft errors can delay computations of the system, the correctness of computations is still guaranteed. Actually the DD scheme is a sufficient construction to get soft-error tolerance, and for optimization, some of double-checking C-elements in DD circuits can be omitted without breaking soft-error tolerance.

### C. Dual Interlocked Programmable Bits

Generally programmable SRAM bits in FPGA employ two conventional cross-coupled inverters, which consist of 6 transistors, as shown in Figure 3 (a). Although the DD scheme can be adapted to obtain soft-error tolerance of programmable bits, a small modification of Calin’s dual interlocked (DI) 12-transistor memory cell [8] is more efficient and also suitable to the DD scheme.

Four nodes (e.g.,  $c_a, \bar{c}_a, c_b, \bar{c}_b$ ) in a DI programmable bit encode a bit as two pairs of complementary values (i.e., 0101, 1010), as shown in Figure 3 (b). The logic state of each node is controlled by two complementary adjacent nodes, and one of the adjacent nodes always keeps a correct value to restore corrupt data. For example, an error at  $c_a$  in 0101 state causes  $\bar{c}_b$  to be an unknown state because both pullup and pulldown transistors are turned on, but  $\bar{c}_a$  still holds 1 that is used to restore  $c_a$ . Other erroneous cases can be similarly analyzed. Each DI memory can tolerate any single error or two errors in pair of non-adjacent nodes (i.e.,  $\{c_a, c_b\}, \{\bar{c}_a, \bar{c}_b\}$ ).

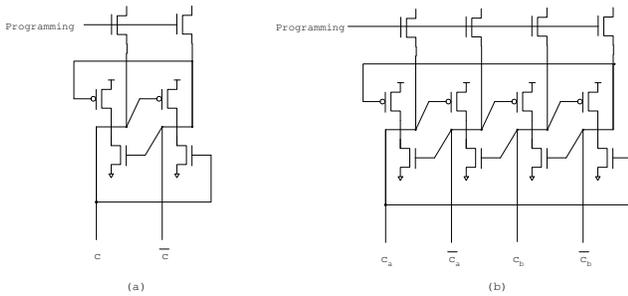


Fig. 3. (a) Programmable Bit (b) Dual Interlocked Programmable Bit

### D. Soft-error tolerant FPGA Cell

We can construct a FPGA logic cell based on the DD scheme with DI programmable bits. Figure 4 shows the construction of a soft-error tolerant logic cell. Two original cells are interweaved by double-checking C-elements. In the same manner, we can construct soft-tolerant S-boxes and C-boxes. The area of the whole FPGA will be enlarged approximately by a factor of two.

## III. CONCLUSION

When a soft error happens, a QDI FPGA may perform incorrectly or halt. The method of duplicating and double-checking nodes provides a way of protecting QDI systems

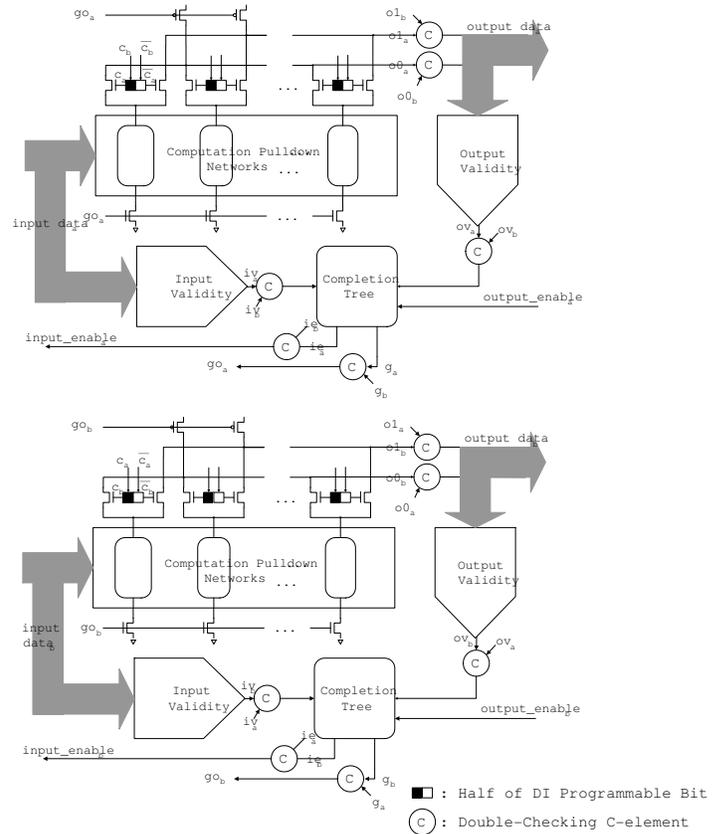


Fig. 4. Soft-error Tolerant FPGA Logic Cell

from soft errors. The stability property of QDI systems permits us to avoid triplicating logic cells. The dual interlocked design is used for programmable bits, and it is more efficient than direct application of the DD scheme to programmable bits. Because the design of soft-error tolerant logic cells is straightforward, we can easily convert existing circuits into soft-error tolerant design and adapt existing synthesis procedures to the duplicated cells.

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